

Symmetry Breaking

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Outline

- What
 - What is *symmetry*?
- Why
 - Why is *symmetry* a problem?
- How
 - How do we deal with *symmetry*?

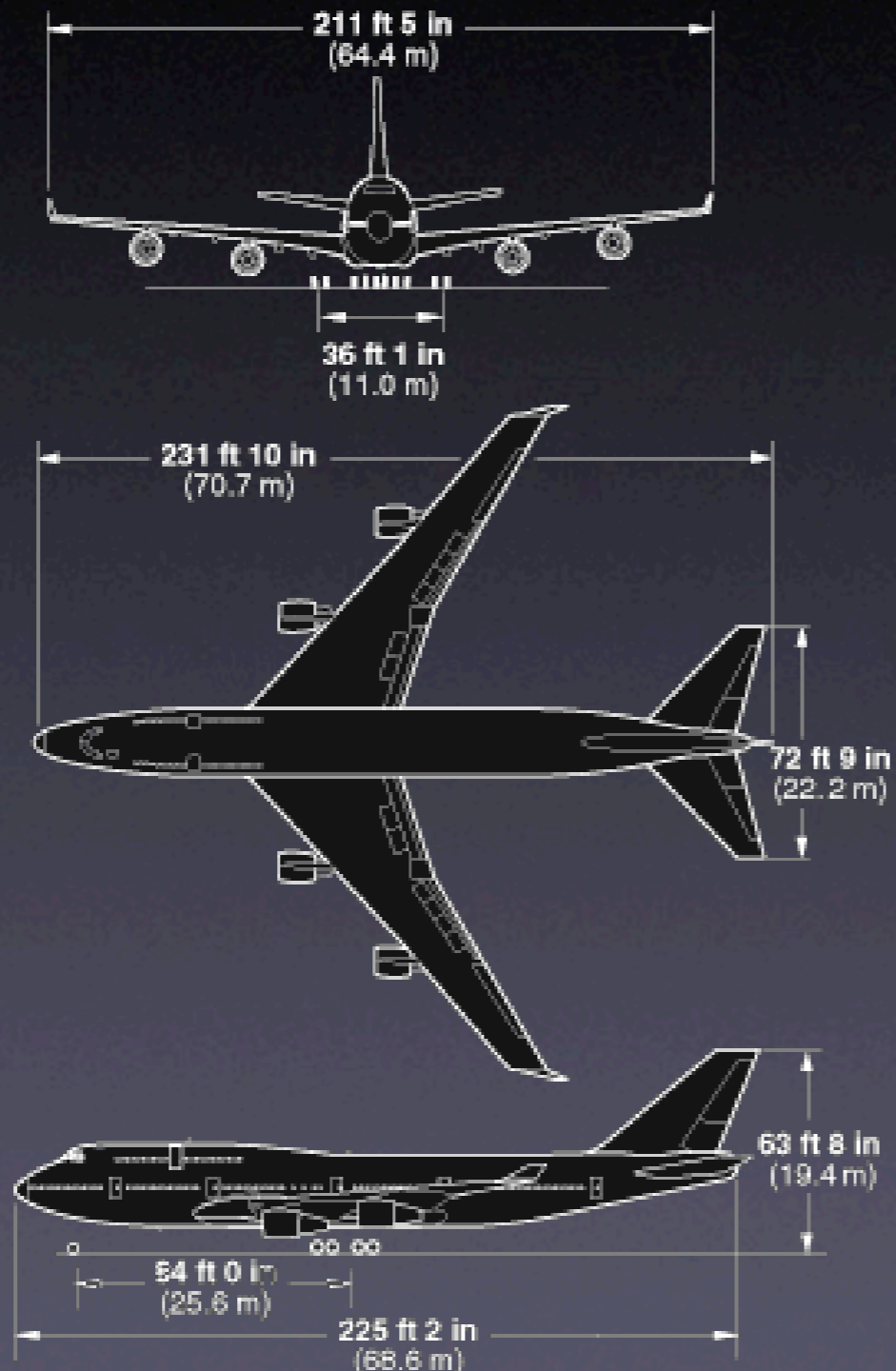
Apology

- Symmetry in constraint programming
- But similar ideas will apply to other domains:
 - Combinatorial optimization
 - Planning
 - Search
 - ...

Active research area

- SymCon'01 workshop, Cyprus 2001
- SymCon'02 workshop, Ithaca 2002
- SymCon'03 workshop, Kinsale 2003
- SymCon'04 workshop, Toronto 2004
- SymCon'05 workshop, Sitges 2005
- SymCon'06 workshop, Nantes 2006
- 1st International Symmetry Conference, Edinburgh 2007

Symmetry



- Within objects
- Design an airplane
- Boeing 747 has reflection symmetry

Symmetry

- Between objects
 - Scheduling problem
 - Fleet of identical 747's



Graph colouring

- Variable for each county
 - Italy, France, ...
 - Values are colours
- Constraints
 - $\text{Italy} \neq \text{Switzerland}$,
 $\text{Italy} \neq \text{France}$, ..



Graph colouring

- Proper colouring
 - Italy=green
 - France=blue
 - Spain=red



Graph colouring

- Symmetric colouring
 - Italy=blue
 - France=red
 - Spain=green



Graph colouring

- Symmetric colouring
- If there are m colours
- $m!$ symmetric solutions



Peaceable armies of coexisting queens



- Place 9 queens and 1 king of each colour on chessboard
- No piece to attack another of the opposite colour

Armies of queens

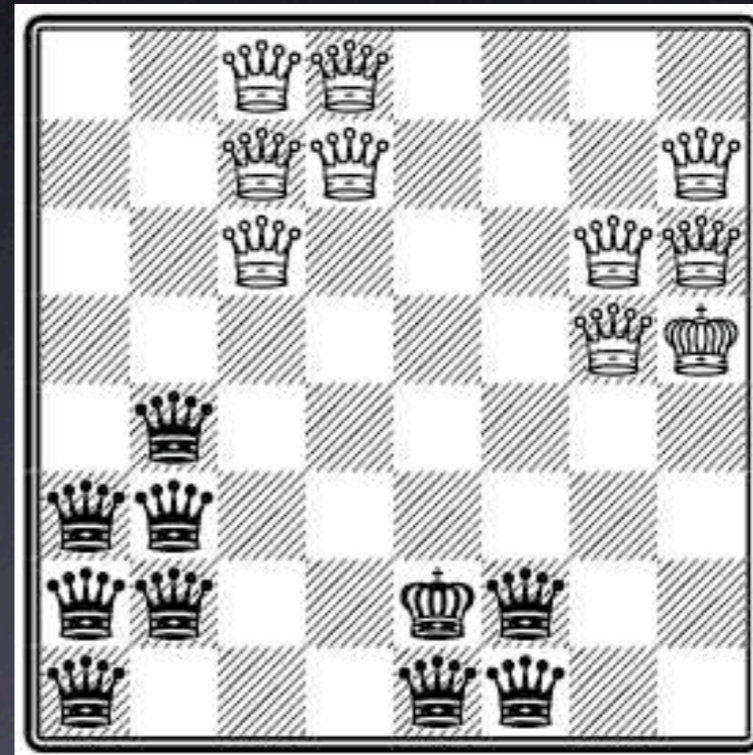
- Set of variables
 - $X[i,j]$ for the square on i th row, j th col
- Set of values
 - {white queen, black queen, empty}

Armies of queens

- Set of constraints
 - $X[i,j]=\text{white queen} \Rightarrow X[i,k] \neq \text{black queen}$
 - $X[i,j]=\text{white queen} \Rightarrow X[k,j] \neq \text{black queen}$
 - $X[i,j]=\text{white queen} \Rightarrow X[i+1,j+1] \neq \text{black queen}$
 - ...

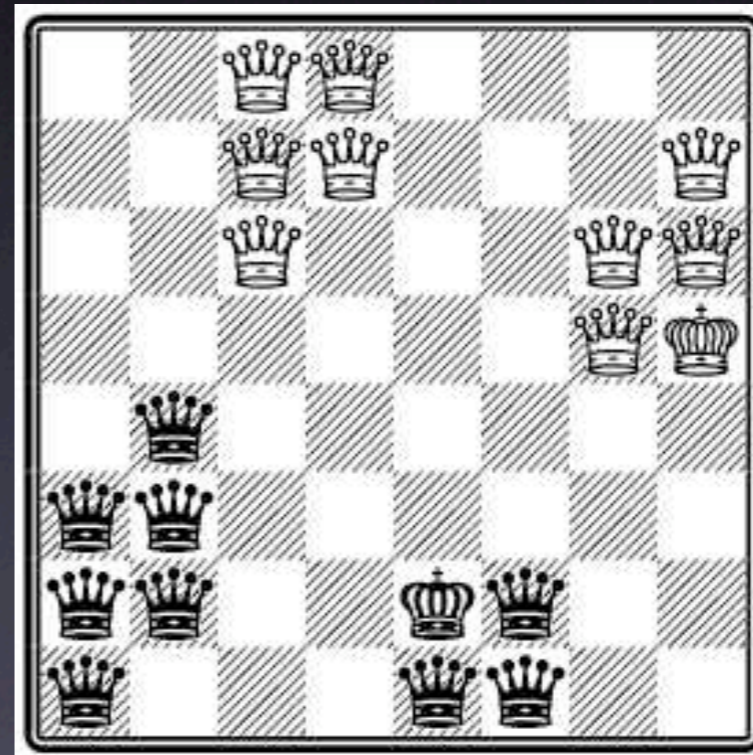
Peaceable armies

- Gives this to a constraint solver
- Here's one solution!



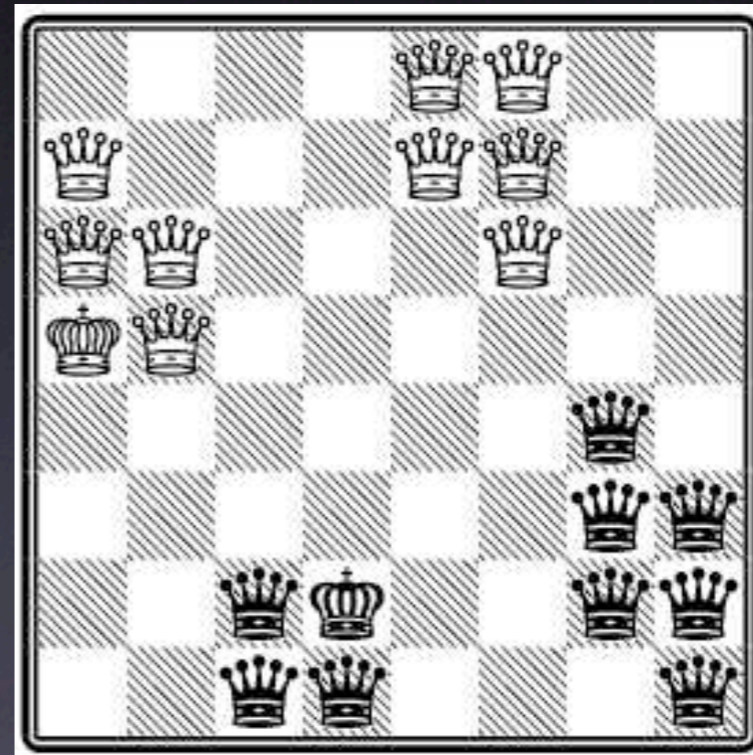
Peaceable armies

- Symmetries of chessboard give other solutions
- horizontal reflection



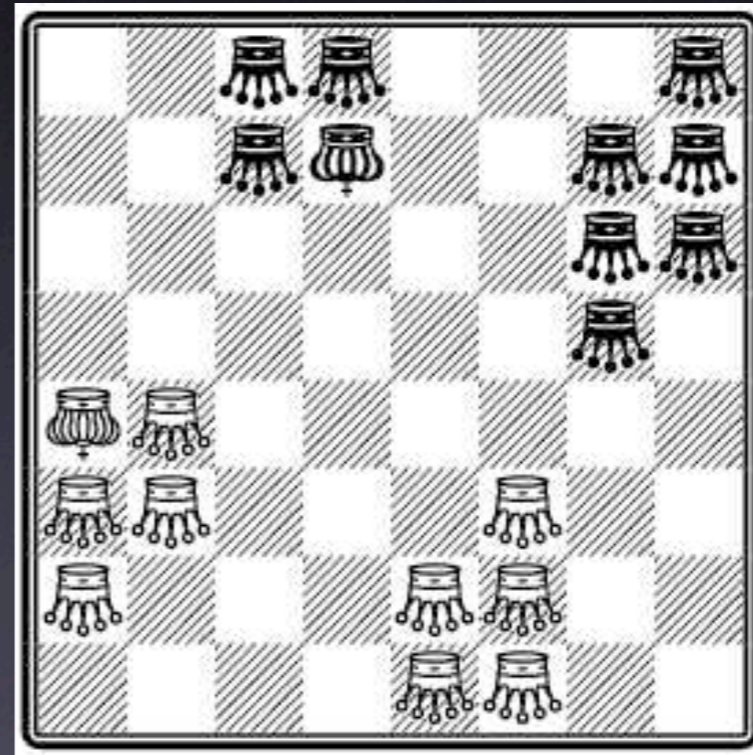
Peaceable armies

- Symmetries of chessboard give other solutions
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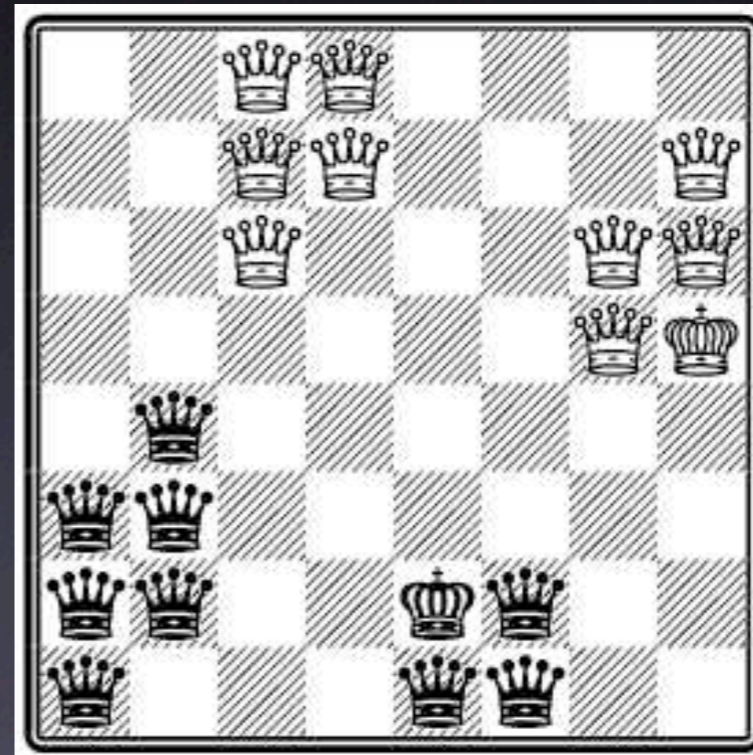
Peaceable armies

- Symmetries of chessboard give other solutions
 - vertical reflection



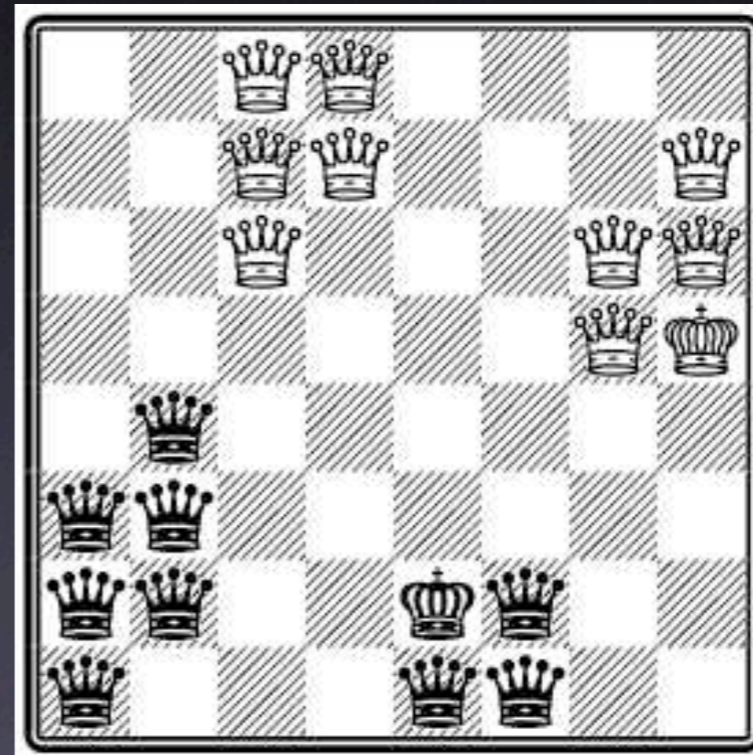
Peaceable armies

- Symmetries of chessboard give other solutions
 - diagonal reflections



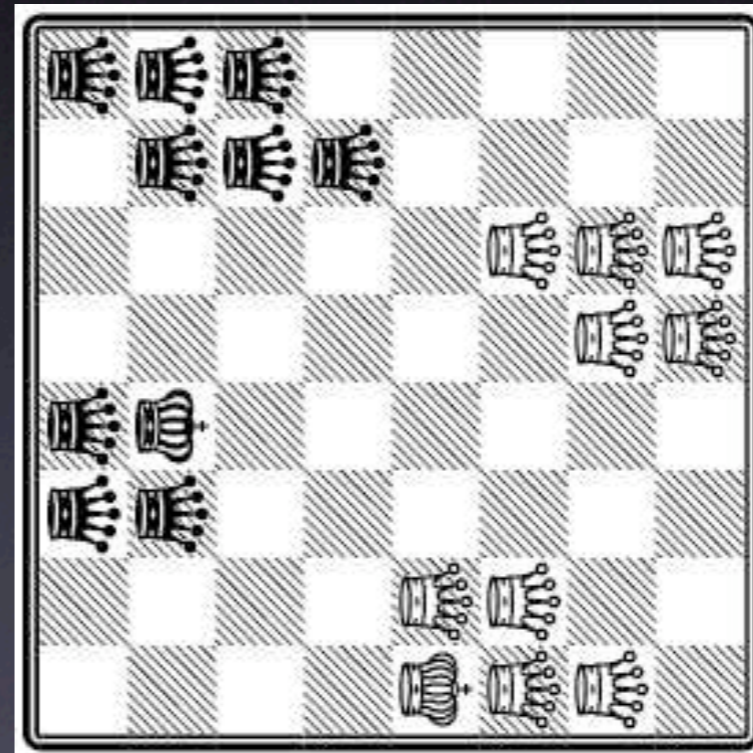
Peaceable armies

- Symmetries of chessboard give other solutions
 - rotation 90 degrees



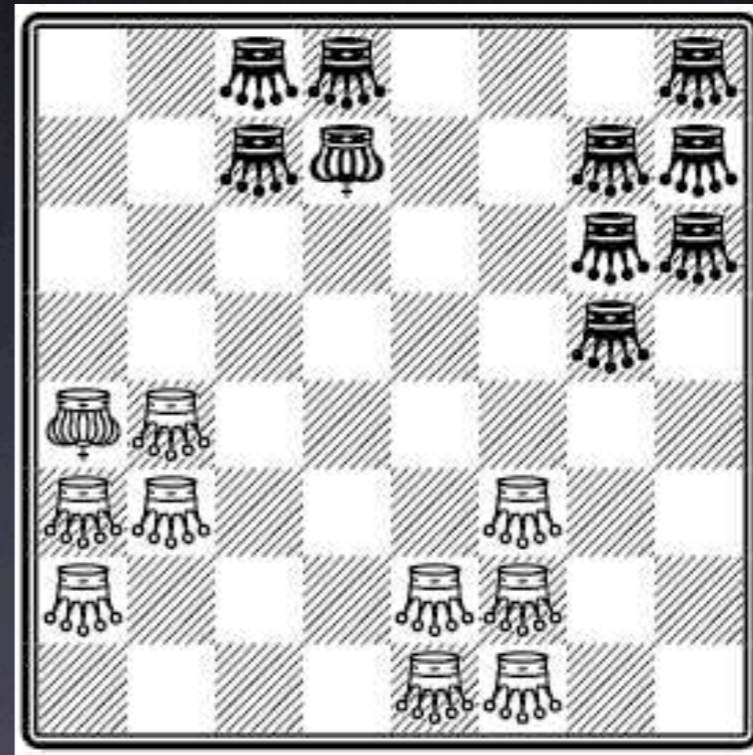
Peaceable armies

- Symmetries of chessboard give other solutions
 - rotation 90 degrees



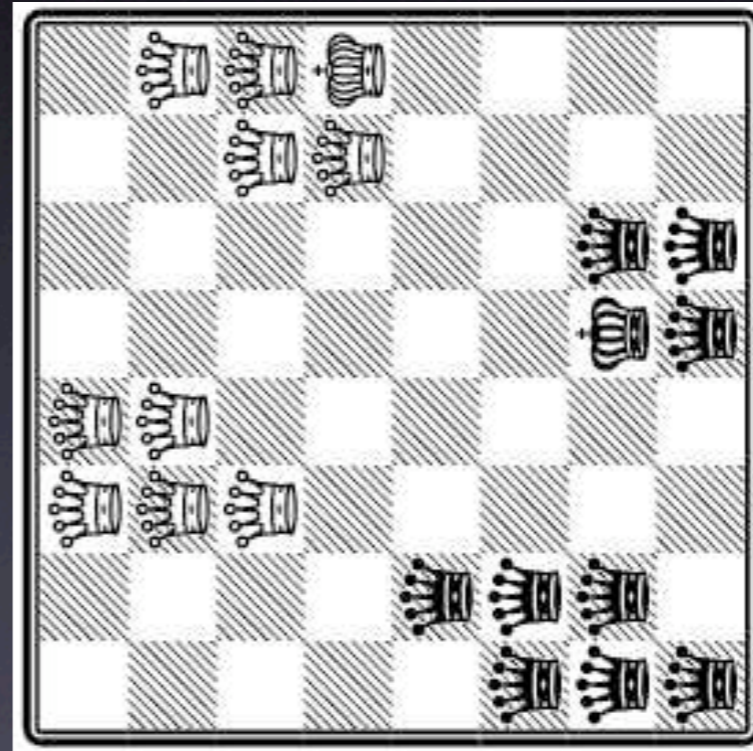
Peaceable armies

- Symmetries of chessboard give other solutions
 - rotation 180 degrees



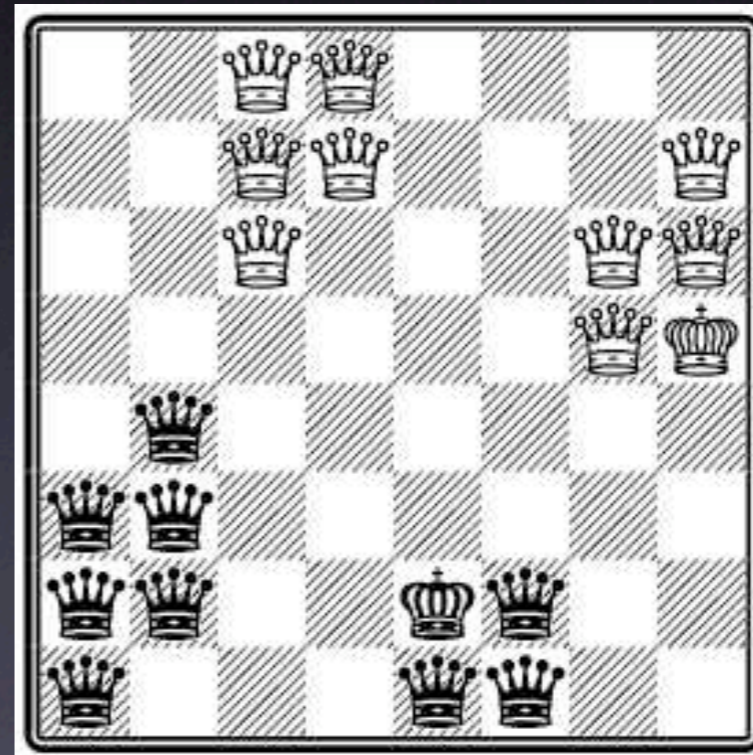
Peaceable armies

- Symmetries of chessboard give other solutions
 - rotation 270 degrees



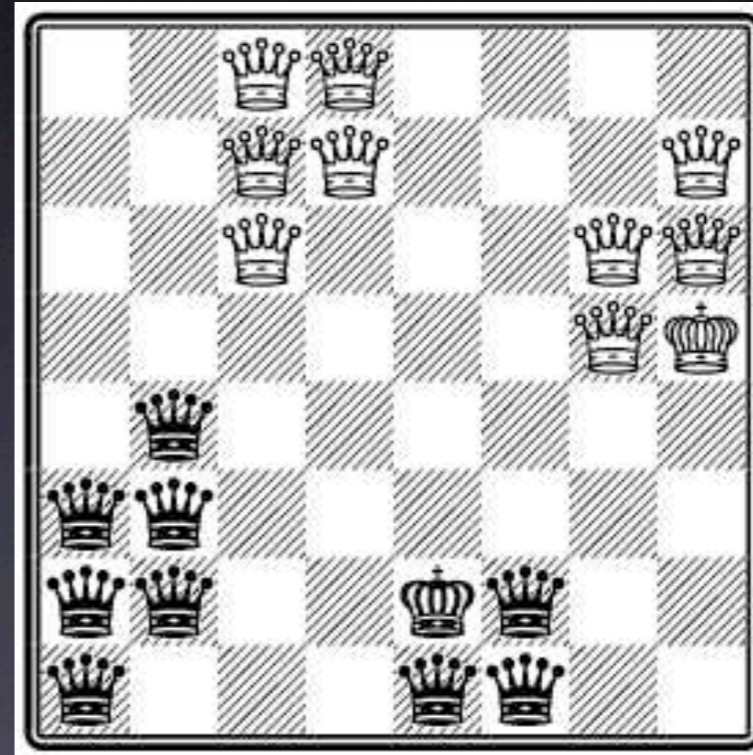
Peaceable armies

- Symmetries of pieces
 - permute any pair of white (or black) queens
 - permute all white pieces with black



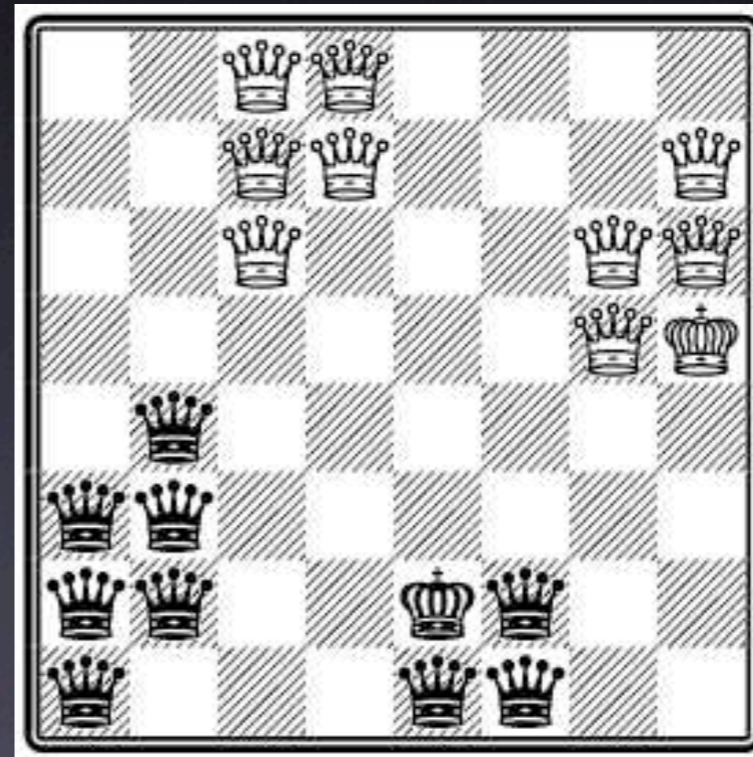
Peaceable armies

- Difficult optimization problem
- Unique solution up to symmetry!
- 2,106,910,310,400 symmetric solutions
- 1/4 US national debt in \$



Peaceable armies

- Difficult optimization problem
- Unique solution up to symmetry!
- 2,106,910,310,400 symmetric solutions
- Don't want to visit symmetric search states



Social golfers



- 32 golfers play once a week in a foursome
- Each week they want to meet 3 different people
- How many weeks can they play?

Social golfers



- 11 weeks is infeasible
- You meet 3 new players each week
- There are only 31 other players
- 10 weeks is possible

Social golfers

- Difficult optimization problem
 - Golfers symmetric
 - Weeks symmetric
 - Order of groups and of foursome irrelevant
 - $32! \times 10! \times 8! \times 4!$ symmetries =
923988455532966699771808443174160957
440000000000 = (mass of Universe in kg)

Social golfers

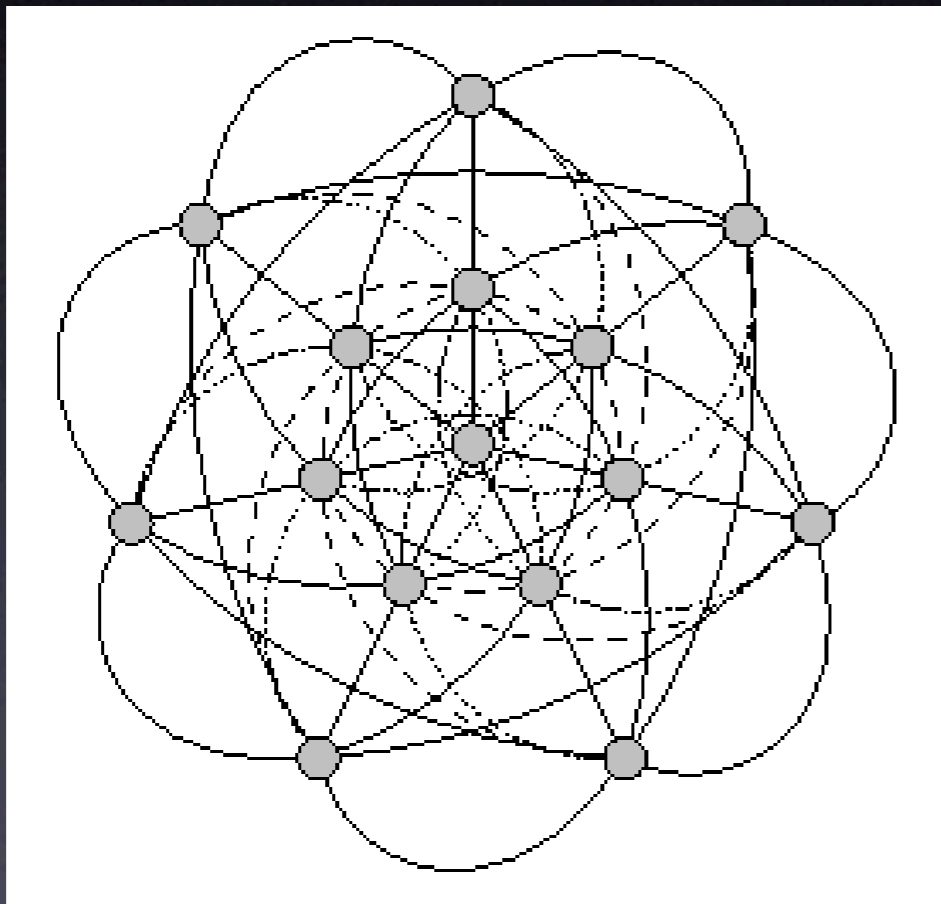
- Simple generalization: (g,s,w) problem
 - g groups
 - groups of size s
 - w weeks

Schoolgirl problem

- Proposed by Rev. Thomas Penyngton Kirkman in the “Ladies and Gentleman’s diary” in 1850
- 15 girls walk in 5 groups of 3 each day for a week. How can the girls be arranged so they walk together with different girls?

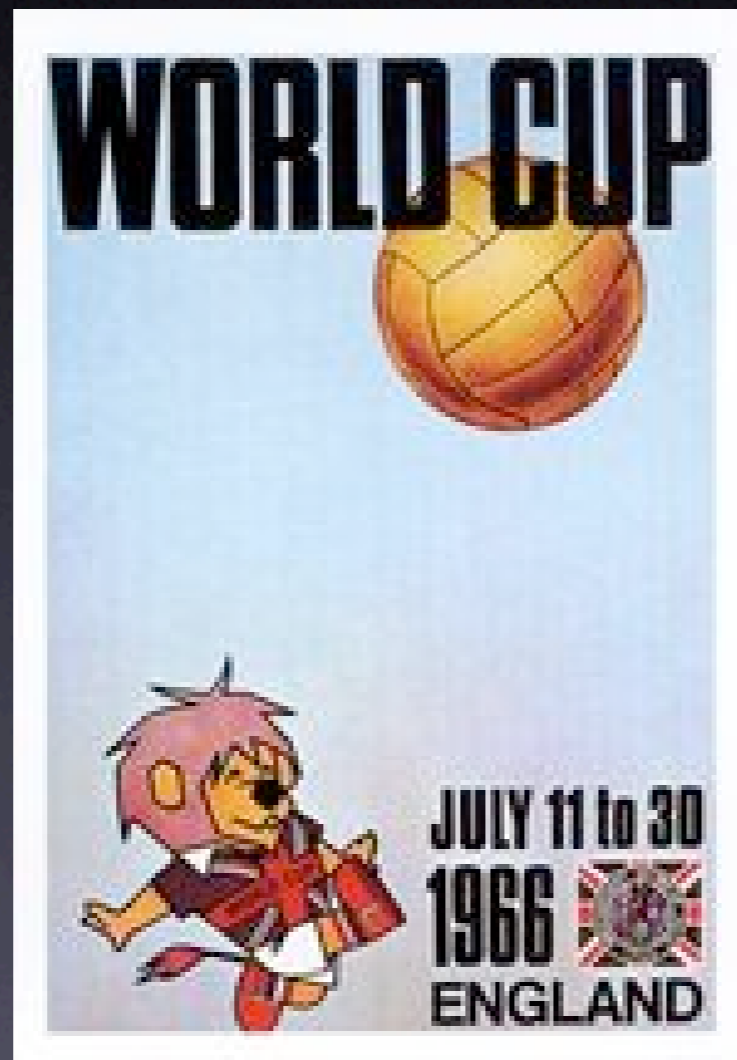


Schoolgirl problem



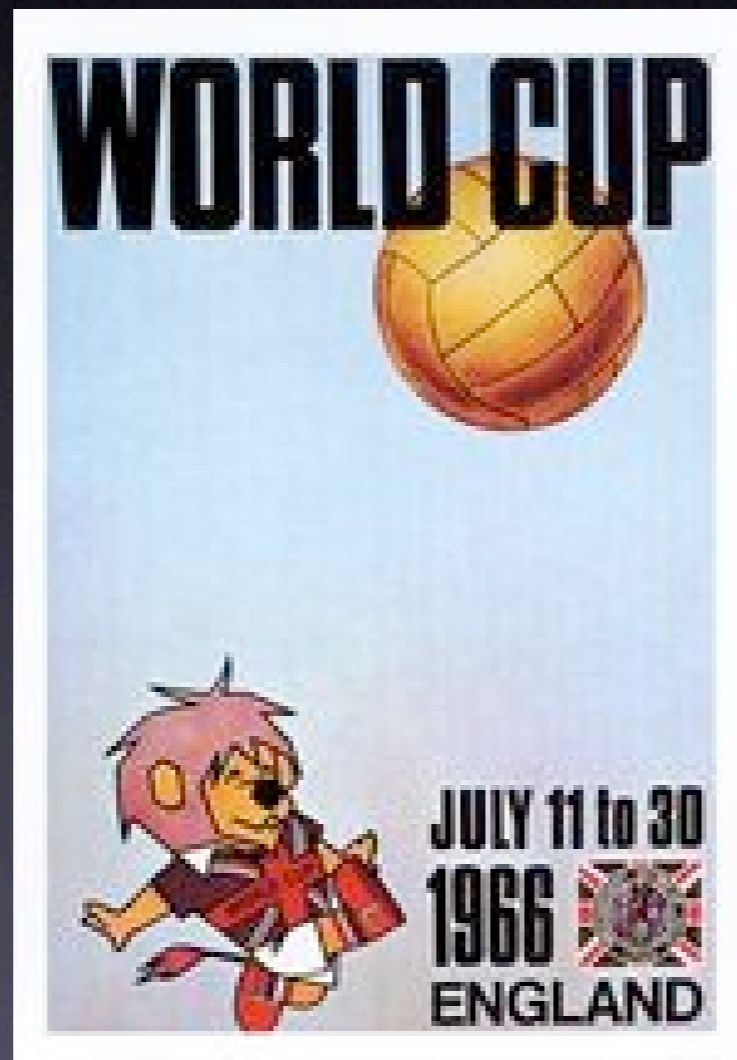
- $(5,3,7)$ problem
- Special type of balanced incomplete block design
- Again lots of symmetry
- girls, groups, days, ...

Tournament scheduling



- Imagine scheduling an event like 1st round of World Cup
 - Suppose 4 venues
 - 8 teams
 - 7 matches (or rounds)

Tournament scheduling



- Often lots of other constraints
- “Home” v “Away” matches
- TV rights
- ..

Tournament scheduling

- Set of variables
 - $\text{Match}[i,j]$ is match played in venue i on round j
- Set of values
 - $\{A \text{v} B, A \text{v} C, A \text{v} D, \dots\}$

Tournament scheduling

- Set of constraints
 - Each team plays once in each round
 - Each team plays every other team
 - ...

Tournament scheduling

- Again lots of symmetry
 - Venues
 - Teams
 - Rounds
- $4! \times 8! \times 7! = 4,877,107,200$



Symmetry

- Scheduling
 - Identical machines, orders
- Rostering
 - Equally skilled workers
- Vehicle routing
 - Identical trucks



Symmetry

- Define in terms of *bijection* on assignments
- Bijection is mapping $\sigma:A \rightarrow B$ that is:
 - Injective: $\sigma(x)=\sigma(y) \Rightarrow x=y$
 - Surjective (onto): $\forall b \in B \exists a \in A . \sigma(a)=b$
 - Also known as *permutation* when $A=B$

Symmetry

- Bijection $\sigma:A \rightarrow A$
 - $A = \{ \langle \text{Italy,red} \rangle, \langle \text{Italy,blue} \rangle, \langle \text{France,red} \rangle, \langle \text{France,blue} \rangle, \dots \}$
 - $\sigma(\langle \text{Italy,red} \rangle) = \langle \text{Italy,blue} \rangle$
 - $\sigma(\langle \text{Italy,blue} \rangle) = \langle \text{Italy,red} \rangle$
 - $\sigma(\langle \text{France,red} \rangle) = \langle \text{France,blue} \rangle$
 - ..

Symmetry in CP

- Solution symmetry
 - Bijection on assignments that preserves solutions (and non-solutions)
- Constraint symmetry
 - Bijection on assignments that preserves constraints

Symmetry in CP

- Solution symmetry
 - $\text{even}(X1+X2), \text{even}(X2+X3)$
 - consider $\sigma(\langle X2, * \rangle) = \langle X3, * \rangle$
- Constraint symmetry
 - $\text{even}(X1+X2), \text{even}(X2+X3), \text{even}(X1+X3)$

Symmetry in CP

- Solution symmetry
 - constraint symmetries \subset solution symmetries
- Constraint symmetry
 - Often the type of symmetries found automatically (using graph isomorphism)

Rotation symmetry

- Symmetry is bijection, σ on assignments that preserves solutions
 - 90 degree rotation
 - $X[1,1]=\text{white queen}, X[2,3]=\text{black queen}$
.. $\Rightarrow X[1,8]=\text{white queen}, X[3,7]=\text{black queen}$..

Permutation symmetry

- Symmetry is bijection, σ on assignments that preserves solutions
 - Permute venues
 - $\text{Match}[1,1]=A \vee B, \text{Match}[2,1]=C \vee D \Rightarrow$
 $\text{Match}[2,1]=A \vee B, \text{Match}[1,1]=C \vee D$..

Permutation symmetry

- Symmetry is bijection, σ on assignments that preserves solutions
 - Permute teams
 - $\text{Match}[1,1]=A \vee B, \text{Match}[2,1]=C \vee D \Rightarrow$
 $\text{Match}[1,1]=A \vee C, \text{Match}[2,1]=B \vee D$..

Types of symmetry

- Variable symmetry
- Value symmetry
- Variable/value symmetry

Types of symmetry

- Variable symmetry
 - Only variables are changed
 - E.g. rotations or reflections of chessboard
 - $X[1,1] \Rightarrow X[1,8]$, $X[2,3] \Rightarrow X[3,7]$
 - Often represent this by permutation of variable indices
 - $(Z[1], Z[2], \dots) \Rightarrow (Z[\sigma(1)], Z[\sigma(2)], \dots)$

Types of symmetry

- Value symmetry
 - Only values are changed
 - E.g. white queen \Rightarrow black queen
 - E.g. $A \vee B \Rightarrow A \vee C$, $C \vee D \Rightarrow B \vee D$
 - In general, $(Z[1], Z[2], \dots) \Rightarrow (\sigma(Z[1]), \sigma(Z[2]), \dots)$

Types of symmetry

- Symmetry can act on both variables and values simultaneously
- E.g. 90 degree rotation of 8-Queens problem
- $\text{Row}[1]=\text{col}2 \Rightarrow \text{Row}[2]=\text{col}8, ..$

Set of symmetries

- Set of symmetries forms a group
- Symmetry breaking exploits group theory
 - generators
 - stabilizers
 - ...

Groups

- Group is set of objects S , and a binary operation •
 - closure: $\forall a, b \in S . a \bullet b \in S$
 - associativity: $\forall a, b, c \in S . (a \bullet b) \bullet c = a \bullet (b \bullet c)$
 - identity: $\exists e \in S \forall a \in S . e \bullet a = a \bullet e = a$
 - inverse: $\forall a \in S \exists b \in S . a \bullet b = b \bullet a = e$

Examples of groups

- C_2 :
 - $\{e, s\}$ where $s \circ s = e$
- C_4 :
 - $\{e, s, s^2, s^3\}$ where $s \circ s = s^2$, $s^2 \circ s = s^3$, $s^3 \circ s = e$

Examples of groups

- C_2 :
 - $\{\text{id}, \text{reflect}\}$ where $\text{reflect} \cdot \text{reflect} = \text{id}$
- C_4 :
 - $\{\text{id}, r_{90}, r_{180}, r_{270}\}$ where $r_{90} \cdot r_{90} = r_{180}$,
 $r_{180} \cdot r_{90} = r_{270}$, $r_{270} \cdot r_{90} = \text{id}$

Example of groups

- Group is set of symmetries S , and a binary operation \cdot which is composition
- closure: since solution/constraints preserved
- associativity: composition is associative
- identity: leave assignments unchanged
- inverse: invert bijection

Permutation group

- Consider permutations of the set $\{1,2,3\}$
 - $e = \text{identity}$, so $e(1)=1, e(2)=2, e(3)=3$
 - $a = (12)$, so $a(1)=2, a(2)=1, a(3)=3$
 - $b = (23)$, so $b(1)=1, b(2)=3, b(3)=2$
- $S_3 = \{e, a, b, ab, ba, aba\}$ forms a group under composition of permutations

Permutation group

- Consider value symmetry in 3 colouring from the set $\{r,g,b\}$
 - $e = \text{identity}$
 - $a = (r\ g)$
 - $b = (g\ b)$
- $S_3 = \{e,a,b,ab,ba,aba\}$ gives the 6 possible permutations of the 3 colours

Group theory

- Generators
 - $\{e, a, b\}$ generates $S_3 = \{e, a, b, ab, ba, aba\}$
 - $a=(1\ 2), b=(2\ 3)$
- Not necessarily unique
 - $\{e, a, aba\}$ also generates S_3
 - $a=(1\ 2), aba=(1\ 3)$

Dealing with *symmetry*

- Don't want to visit *symmetric* search states
 - “Identical” solutions
 - “Identical” failing states
- How do we eliminate these from search?

Reformulation

- Change representation
 - $\text{WhiteQueen}[1]=(1,1)$, $\text{WhiteQueen}[2]=(1,2)$, ..., $\text{BlackQueen}[1]=(5,7)$, ..
 - $X[1,1]=\text{white queen}$, $X[1,2]=\text{white queen}$, ..., $X[5,7]=\text{black queen}$

All interval series

- Order numbers 0 to $n-1$ so that
 - Each difference between neighbouring numbers occurs once
 - E.g. 0 8 1 7 2 6 3 5 4
 - Diff: 8 7 6 5 4 3 2 1
 - What symmetries does this problem have?

All interval series

- Order numbers 0 to $n-1$ so that
 - Each difference between neighbouring numbers occurs once
 - E.g. 0 8 1 7 2 6 3 5 4
 - Diff: 8 7 6 5 4 3 2 1
 - Reversal symmetry: 4 5 3 6 2 7 1 8 0

All interval series

- Order numbers 0 to $n-1$ so that
 - Each difference between neighbouring numbers occurs once
 - E.g. 0 8 1 7 2 6 3 5 4
 - Diff: 8 7 6 5 4 3 2 1
 - Complementation: 8 0 7 1 6 2 5 3 4

Reformulation of AIS

- Cyclic view
 - Order numbers 0 to $n-1$ in a cycle
 - Each difference 1 to $n-1$ occurs
 - E.g. 0 8 1 7 2 6 3 5 4
 - Diffs: 8 7 6 5 4 3 2 1 4
 - What symmetries does this now have?

Reformulation of AIS

- Cyclic view
 - Order numbers 0 to $n-1$ in a cycle
 - Each difference 1 to $n-1$ occurs
 - E.g. 0 8 1 7 2 6 3 5 4
 - Diffs: 8 7 6 5 4 3 2 1 4
 - Reversal symmetry

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Reformulation of AIS

- Cyclic view
 - Order numbers 0 to $n-1$ in a cycle
 - Each difference 1 to $n-1$ occurs
 - E.g. 0 8 1 7 2 6 3 5 4
 - Diffs: 8 7 6 5 4 3 2 1 4
 - Rotation symmetry

Reformulation of AIS

- Cyclic view
 - Order numbers 0 to $n-1$ in a cycle
 - Each difference 1 to $n-1$ occurs
 - E.g. 0 8 1 7 2 6 3 5 4
 - Diffs: 8 7 6 5 4 3 2 1 4
 - Symmetry easily broken: sequence starts
0 $n-1$ 1

Reformulation of AIS

- E.g. 0 8 1 7 2 6 3 5 4
- Diffs: 8 7 6 5 4 3 2 1 4
 - Given solution to cyclic view
 - reverse: 4 5 3 6 2 7 1 8 0
 - complement: 8 0 7 1 6 2 5 3 4
 - both: 4 3 5 2 6 1 7 0 8
 - common diff: 6 3 5 4 0 8 1 7 2 (and its symmetries)

Breaking symmetry

- Add symmetry breaking constraints
 - $\text{Match}[1,1]=A \vee B$
 - $\text{Match}[2,1]=C \vee D$
 - ...

Rehearsal problem

- $X[i]$ = scene rehearsed in i th time slot
 - Actors must arrive before their first scene and stay till their last scene
- Reflection symmetry
 - Can reverse any rehearsal sequence
 - Prevent this with constraint: $X[1] < X[n]$

LEX LEADER

- For variable symmetries, [Crawford et al. KR96] give general method:
 - Pick order on vars: $X[1]$ to $X[n]$
 - For each variable symmetry σ , post LEX LEADER constraint:
 - $(X[1], \dots, X[n]) \leq_{\text{lex}} (X[\sigma(1)], \dots, X[\sigma(n)])$

Lexicographical order

- $(Y_1, Y_2, \dots) \leq_{\text{lex}} (Z_1, Z_2, \dots)$ iff
 - $Y_1 < Z_1$ or
 - $Y_1 = Z_1$ & $(Y_2, \dots) \leq_{\text{lex}} (Z_2, \dots)$
- Order used in dictionaries, etc
 - $(1, 1, 2, 1, 2, 3, 1, \dots) \leq_{\text{lex}} (1, 1, 3, 1, 3, 2, 1, \dots)$
- Linear time propagator [Frisch, Hnich, Kiziltan, Miguel, Walsh CP02]

Rehearsal problem

- $X[i]$ = scene rehearsed in i th time slot
 - Actors must arrive before their first scene and stay till their last scene
- Reflection symmetry
 - Can reverse any rehearsal sequence
 - $(X[1], \dots, X[n]) \leq_{\text{lex}} (X[n], \dots, X[1])$
 - Simplifies to $X[1] < X[n]$

Non-attacking queens

- $X[i,j] \in \{\text{white queen, black queen, empty}\}$
- 90 rotation symmetry
- $(X[1,1], X[1,2], \dots, X[1,8], X[2,1], \dots, X[2,8], \dots) \leq_{\text{lex}} (X[8,1], X[7,1], \dots, X[1,1], X[8,2], \dots, X[1,2], \dots)$

Non-attacking queens

- $X[i,j] \in \{\text{white queen, black queen, empty}\}$
- 180 rotation symmetry
- $(X[1,1], X[1,2], \dots, X[1,8], X[2,1], \dots, X[2,8], \dots) \leq_{\text{lex}}$
 $(X[8,8], X[8,7], \dots, X[8,1], X[7,8], \dots, X[7,1], \dots)$

Non-attacking queens

- $X[i,j] \in \{\text{white queen, black queen, empty}\}$
- 270 rotation symmetry
- $(X[1,1], X[1,2], \dots, X[1,8], X[2,1], \dots, X[2,8], \dots) \leq_{\text{lex}}$
 $(X[1,8], X[2,8], \dots, X[8,8], X[1,7], \dots, X[8,7], \dots)$

Non-attacking queens

- $X[i,j] \in \{\text{white queen, black queen, empty}\}$
- horizontal reflection
 - $(X[1,1], X[1,2], \dots, X[1,8], X[2,1], \dots, X[2,8], \dots) \leq_{\text{lex}} (X[8,1], X[8,2], \dots, X[8,8], X[7,1], \dots, X[7,8], \dots)$

Non-attacking queens

- $X[i,j] \in \{\text{white queen, black queen, empty}\}$
- vertical reflection
 - $(X[1,1], X[1,2], \dots, X[1,8], X[2,1], \dots, X[2,8], \dots) \leq_{\text{lex}} (X[1,8], X[1,7], \dots, X[1,1], X[2,8], \dots, X[2,1], \dots)$

LEX LEADER method

- Three challenges
 - Extend method to work with other types of symmetry (e.g. value symmetries)
 - Deal with exponential number of symmetries
 - Conflict between branching heuristic and symmetry breaking constraints

Variable symmetry

- Bijection σ on vars which maps solutions onto solutions
 - E.g. reflection symmetry:
 $X[1] \rightarrow X[n], X[2] \rightarrow X[n-1], \dots$
- LEX LEADER method
 - E.g. $(X[1], \dots, X[n]) \leq_{\text{lex}} (X[n], \dots, X[1])$

Value symmetry

- Bijection ϑ on values which maps solutions onto solutions
- E.g. suppose two scenes have same actors, then can permute these two scenes (=values) in any rehearsal
- LEX LEADER method
 - $(X[1], \dots, X[n]) \leq_{\text{lex}} (\vartheta(X[1]), \dots, \vartheta(X[n]))$

Value symmetry

- Puget's propagator
 - Construct symmetric assignment:
 - E.g. $\text{Element}(X[i], [\vartheta(1), \dots, \vartheta(m)], Y[i])$
 - Lex ordering result
 - $(X[1], \dots, X[n]) \leq_{\text{lex}} (Y[1], \dots, Y[n])$
 - But does not achieve GAC!

Value symmetry

- Linear time GAC propagator
 - $X[1] \ X[2] \ .. \ X[n] \leq_{lex}$
 $\vartheta(X[1]) \ \vartheta(X[2]) \ .. \ \vartheta(X[n])$
 $B[1]=0 \ B[2] \ .. \ B[n] \ B[n+1]$
 - Post $C(X[i], B[i], B[i+1])$ where
 - $B[i]=B[i+1]=0$ and $X[i]=\vartheta(X[i])$, or
 - $B[i]=0, B[i]=1$ and $X[i] < \vartheta(X[i])$, or
 - $B[i]=B[i+1]=1$
 - Example: $X[1] \in \{1,2\}, X[2]=2, \sigma(1)=2, \sigma(2)=1$

Var and value symmetry

- Bijection σ on vars, and bijection ϑ on values that maps solutions to solutions
 - E.g. reversal of rehearsal (var symmetry) and permutation of scenes (val symmetry)
- LEX LEADER method
 - $(X[1], \dots, X[n]) \leq_{\text{lex}} (\vartheta(X[\sigma(1)]), \dots, \vartheta(X[\sigma(n)]))$

Var/value symmetry

- Symmetries may act simultaneously on vars and values
 - Cannot be decomposed into bijection on vars, and bijection on values
 - E.g. in n queens problem, rotate 90°:
 $X[i]=j \rightarrow X[j]=n-i+1$
- Bijection on (vars, values)
 - E.g. $\sigma(i,j) = j, n-i+1$

Var/value symmetry

- Not all (partial) assignments map onto proper (partial) assignments
 - E.g. $X[1]=1, X[2]=1 \dots \rightarrow X[1]=n, X[1]=n-1 \dots$
- LEX leader method
 - Admissible($[X[1], \dots, X[n]]$) &
 $(X[1], \dots, X[n]) \leq_{\text{lex}} \sigma(X[1], \dots, X[n])$

Lots of symmetries

- LEX LEADER method posts one constraint per symmetry
- Can be exponential number of symmetries
- E.g. m indistinguishable values gives $m!$ value symmetries
- How can we deal efficiently and effectively with such situations?

Modifying search

- Avoid visiting symmetric states
 - SBDS (symmetry breaking during search)
 - SBDD (symmetry breaking by dominance detection)
 - GE-trees (group equivalence trees)

Modifying search

- Symmetry Breaking During Search
 - add a constraint at each node to rule out symmetric equivalents in the future
- Symmetry Breaking by Dominance Detection
 - check each node before entering it, to make sure you have not been to an equivalent in the past

SBDS

- Branch
 - Given assignments so far $X[1]=a[1], \dots, X[k-1]=a[k-1]$
 - Try $X[k]=b$

SBDS

- Branch
 - Given assignments so far $X[1]=a[1], \dots, X[k-1]=a[k-1]$
 - Try $X[k]=b$, if this fails
 - Post $X[k] \neq b$

SBDS

- Branch
 - Given assignments so far $X[1]=a[1], \dots, X[k-1]=a[k-1]$
 - Try $X[k]=b$, if this fails
 - Post $X[k] \neq b$, and don't visit a symmetric state to the last branch

SBDS

- Branch
 - Given assignments so far $X[1]=a[1], \dots, X[k-1]=a[k-1]$
 - Try $X[k]=b$, if this fails
 - Post $X[k] \neq b$, if $\sigma(X[1]=a[1]) \& \dots \sigma(X[k-1]=a[k-1])$ then $\neg \sigma(X[k]=b)$

SBDS

- E.g. reflection symmetry
- Given assignments so far $X[1]=a[1], \dots, X[k-1]=a[k-1]$
- Try $X[k]=b$, if this fails
- Post $X[k] \neq b$, if $X[n]=a[1] \ \& \ \dots \ X[n-k+2]=a[k]$ then $X[n-k+1] \neq b$

SBDS

- +ve
 - Does not conflict with branching heuristics
- -ve
 - Need to post symmetry breaking constraint for each symmetry
 - In general, may be exponential number of symmetries

SBDD

- Fahle, Schamberger, Sellmann, 2001
- Focacci, Milano, 2001
- prefigured by Brown, Finkelstein, Purdom, 1988
- do not search a node if you have searched its equivalent before
- check before entering a node

SBDD

- +ve
 - Does not conflict with branching heuristic
- -ve
 - Need to code dominance detection
 - Only “forward checking”
 - Can take exponential time on problems static methods solve without search

Special cases

- Value symmetry
 - Interchangeable values
- Variable symmetry
 - Row and column symmetry

Interchangeable values

- Often we have some (sub)set of values which can be freely interchanged
 - {golfer 1, golfer 2, ...}
 - {white queen, black queen}
- Given solution, we can uniformly swap values

Interchangeable values

- Often we have some (sub)set of values which can be freely interchanged
 - {golfer 1, golfer 2, ...}
 - {white queen, black queen}
- If there are m values, $m!$ symmetries
 - But we can deal with them efficiently and effectively!

Interchangeable variables

- Often we have some (sub)set of variables which can be freely interchanged
 - Queen[1]=(1,2), Queen[2]=(4,3), ..
- Easy to break this symmetry!
 - Order variables, Queen[1] < Queen[2] < ..

Interchangeable vars and values

- Sometimes we can have both interchangeable variables and values
 - Consider graph colouring
 - Node1 = red, Node2 = blue, ..
 - Suppose Node1 and Node2 have the same neighbours

Interchangeable vars and values

- Sometimes we can have both interchangeable variables and values
 - Consider pigeonhole problem
 - Hole1 = pigeon1, Hole2 = pigeon3, ..
 - Holes and pigeons all interchangeable

Row and col symmetries

- Many problems can be modelled with matrix of decision variables
- Combinatorial problems like BIBD
- Rows and cols can be freely permuted

```
0 1 1 0 0 1 0
1 0 1 0 1 0 0
0 0 1 1 0 0 1
1 1 0 0 0 0 1
0 0 0 0 1 1 1
1 0 0 1 0 1 0
0 1 0 1 1 0 0
```

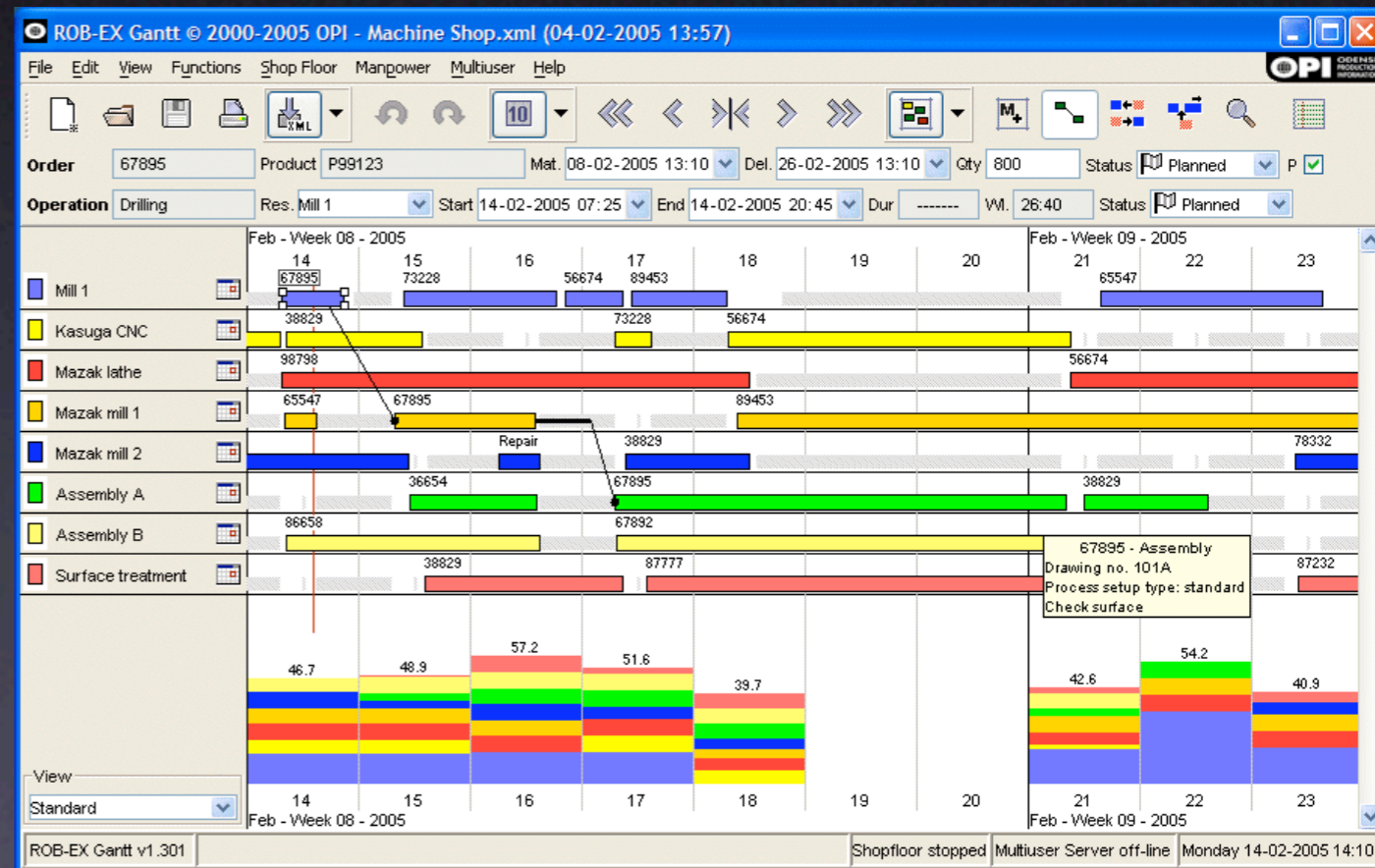

Row and col symmetries

- Many problems can be modelled with matrix of decision variables
- Scheduling problems like social golfer
- Group $[i,j]$ are golfers playing in i th group on week j
- Rows and cols can be freely permuted



Row and col symmetries

- Many problems can be modelled with matrix of decision variables
- Production planning problems
- $\text{Order}[i,j,k]=1$ iff order i goes on machine j in shift k
- Rows and cols can be (partially) permuted



Row and col symmetries

- If we have a n by m matrix of decision variables
 - $m!n!$ row and col symmetries
- However, as we shall see later, efficient and effective means to deal with this large number of symmetries
 - Again uses the LEX constraint!

Outline

- What is symmetry?
 - Bijection on assignments preserving solutions/
constraints
 - Variable and value symmetry
 - Two important special cases
 - Interchangeable values
 - Row and col symmetry

Outline

- Why is symmetry a problem?
 - Increases size of search space!
- How do we deal with symmetry?
 - Reformulate problem
 - Add constraints
 - LEX LEADER method
 - Modify search
 - SBDS, SBDD, GE-tree

Conclusions

- Symmetry occurs in many problems
 - We must deal with it or face a combinatorial explosion!
- We have some generic methods (for small numbers of symmetries)
 - In special cases, we can break all symmetries

Questions?

