

# Complexity & Symmetry

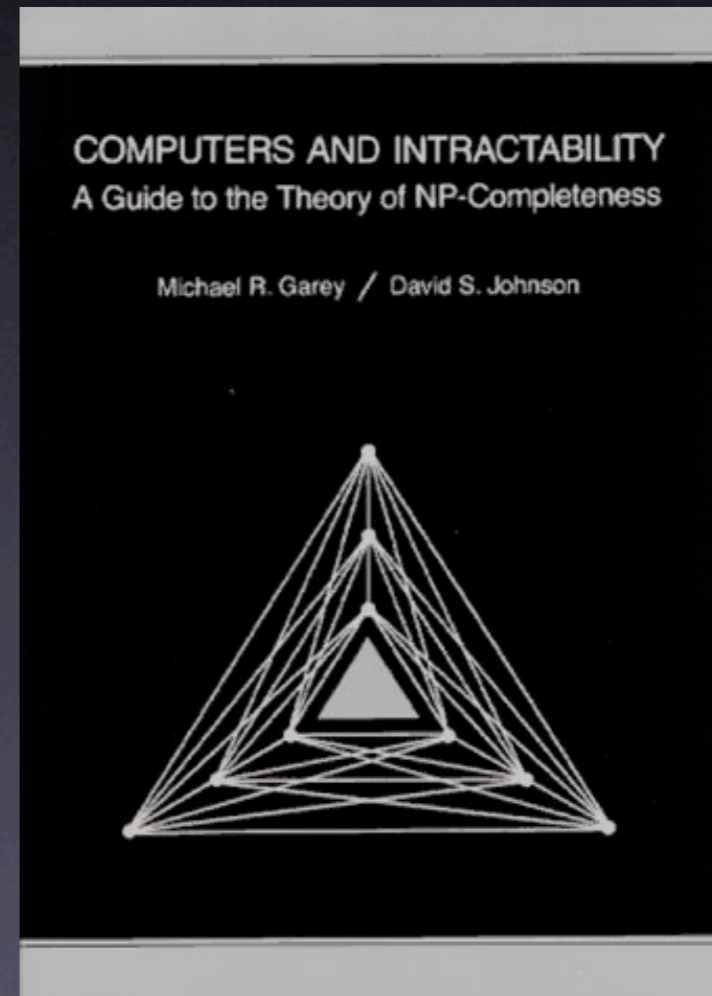
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# Global constraints

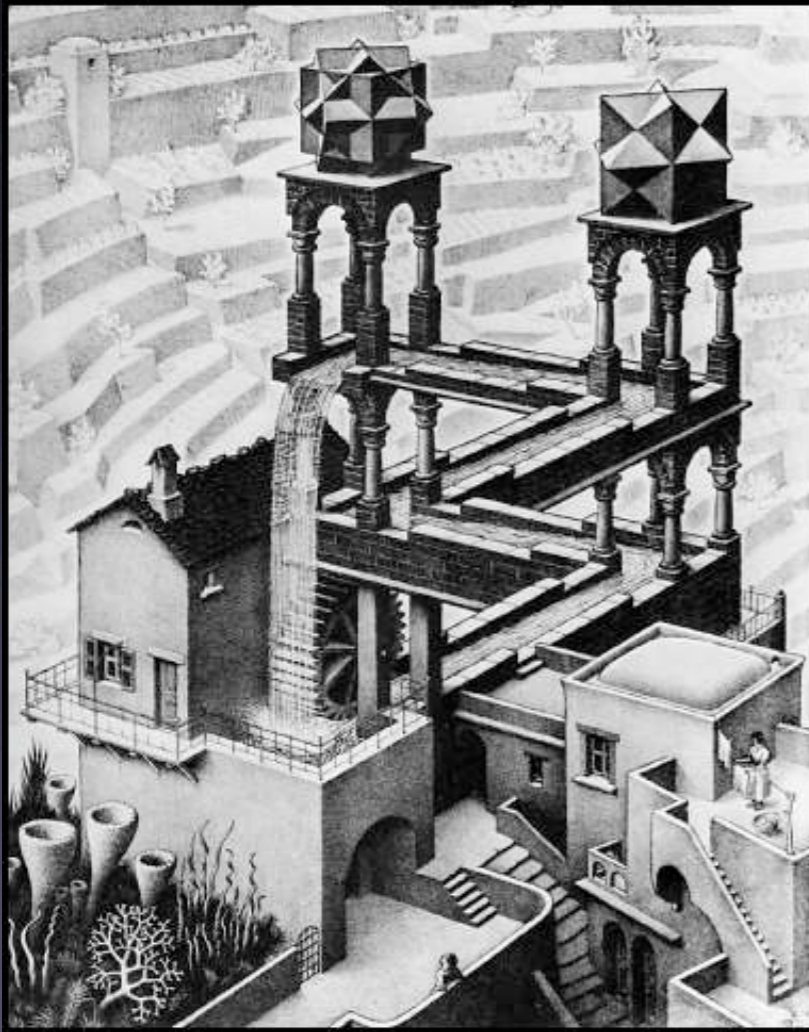
- Capture common patterns
  - $\text{Alldifferent}(X_1, \dots, X_n)$
  - $\text{Nvalues}(N, X_1, \dots, X_n)$
  - $\text{Lex}([X_1, \dots, X_n], [Y_1, \dots, Y_n])$
- Efficient & effective specialized propagators
  - Prune parts of search tree

# Our hammer

- Use basic tools of computational complexity to study limits of
  - reasoning with global constraints
  - global constraints for breaking symmetry



# Limits of Global Constraints



- Enforce lesser consistency
- Constraints cannot be combined
- Constraints cannot be generalized
- Decomposition will hurt pruning

# In general

- Consider (generalized) arc-consistency
  - Every value for every variable can be extended to satisfy the constraint
  - That is, every value has *support*
- Similar results for other local consistencies
  - Bounds consistency for integer variables
  - Bounds consistency for set variables
  - ...

# In general

- Global constraints are intractable
  - ACSupport? is NP-complete
    - Does this value have support?
- Consider  $C(X_1, \dots, X_n)$ 
  - Where  $X_i=j$  implies  $X_j=\text{true}$ ,  $X_i=-j$  implies  $X_j=\text{false}$
  - SAT in  $k$  vars,  $j$  clauses  $\rightarrow C(X_1, \dots, X_{j+k})$ 
    - $X_1$  to  $X_k \in \{\text{true}, \text{false}\}$
    - $i$ th clause is  $x_1 \vee -x_3 \vee x_5 \rightarrow X_{k+i} \in \{1, -3, 5\}$
    - Consider reduction of:  $\{x_1, -x_1 \vee x_2\}$

# In general

- Global constraints are intractable
  - ACSupport? is NP-complete
    - Does this value have support?
  - MaxAC? is DP-complete
    - Are these domains the maximal arc-consistent domains?
    - DP is NP  $\cup$  coNP
      - Answers NP question: are these domains AC? Yes!
      - Answers coNP questions: is any smaller domain AC? No!

# In general

- Global constraints are intractable
  - ACSupport? is NP-complete
  - MaxAC? is DP-complete
- 
- Even some specific constraints proposed in the past are intractable
  - NValues( $N, X_1, \dots, X_n$ )
  - AtMost1( $S_1, \dots, S_n$ )
  - ...



# NValues

- $NValues(N, X_1, \dots, X_n)$ 
  - $N$  values used in  $X_1, \dots, X_n$
  - Useful for resource allocation
- Simple reduction of SAT to NValues
  - SAT problem in  $k$  vars,  $j$  clauses
  - $X_i = \{i, -i\}$  for  $1 \leq i \leq k$
  - $X_{k+i} = \{1, -3, 5\}$  if  $i$ -th clause is:  $(x_1 \vee -x_3 \vee x_5)$
  - $N = \{n\}$
  - Consider reduction of:  $\{x_1, -x_1 \vee x_2\}$

# NValues

- $NValues(N, X_1, \dots, X_n)$ 
  - $N$  values used in  $X_1, \dots, X_n$
  - Useful for resource allocation
- Simple reduction of SAT to NValues
  - Finding support (and hence enforcing arc-consistency) is NP-hard
  - Look to enforce *lesser* level of local consistency like bound consistency

# Composing constraints

- Take two tractable constraints
  - E.g.  $\text{Disjoint}(S_1, \dots, S_n)$  and  $\text{FixedCard}(S_1, \dots, S_n)$
- Could we combine them into one *bigger* global constraint?
  - E.g.  $\text{FixedCardDisjoint}(S_1, \dots, S_n)$
  - No, NP-hard to propagate!

# GCC

- Take a tractable constraint
  - E.g.  $\text{GCC}([X_1, \dots, X_n], [l_1, \dots, l_m], [u_1, \dots, u_m])$
  - Value  $j$  occurs between  $l_j$  and  $o_j$  times in  $X_1, \dots, X_n$
- Generalize some constants to variables
  - E.g.  $\text{GCC}([X_1, \dots, X_n], [O_1, \dots, O_m])$
  - Value  $j$  occurs  $O_j$  times in  $X_1, \dots, X_n$
- NP-hard to make generalized arc-consistent
  - [Claude-Guy Quimper 2003]

# GCC

- Reduction of 1 in 3 SAT on +ve clauses to GCC
- If  $i$ th clause is  $x_1 \vee x_3 \vee x_5$  then  $X_i \in \{1, 3, 5\}$
- $O_j \in \{0, k\}$  where  $k$  is number of occurrences of  $x_j$  in clauses
- Consider  $\{x_1 \vee x_2 \vee x_4, x_2 \vee x_3 \vee x_4, x_1 \vee x_3 \vee x_4\}$

# Decomposing constraints

- Consider a global constraint that is NP-hard to propagate
  - E.g.  $\text{AtMost } 1(S_1, \dots, S_n)$
- Consider a decomposition into smaller constraints
  - $|S_i \cap S_j| \leq 1$  for all  $i < j$
- If it is polynomial to propagate decomposition
  - decomposition must hinder propagation (assuming  $P \neq NP$ )

# Symmetry breaking

- Add (global) constraints to eliminate symmetries
  - E.g. lex order rows, lex order cols
- Can we break *all* row & col symmetry with a single global constraint?
  - Enforcing GAC on such a global constraint is NP-hard

# Conclusions

- Computational complexity is a useful hammer to study global constraints
- Uncovers fundamental limits of reasoning with global constraints
  - Lesser consistency needs to be enforced
  - Decomposition hurts pruning
  - Composition or generalization intractable
  - Symmetry breaking is inherently limited
- ...